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# *Acid Mixes*

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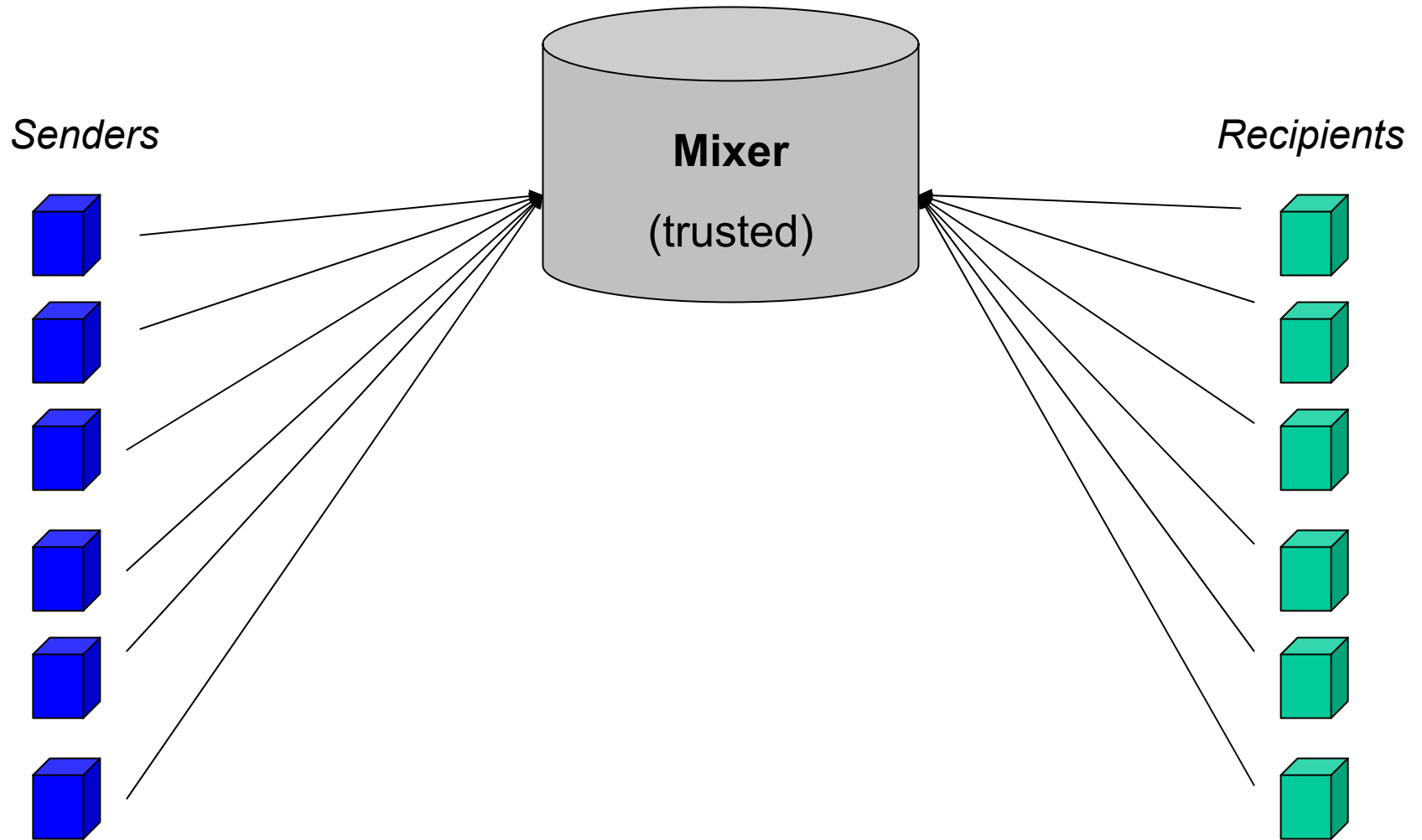
# *What is that?*

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- A variation on mix-net protocols that (attempts to) address *reliability* and *trust* issues while maintaining *anonymity* and preserving *ACID* properties.
- The variation is, itself, a “mix”:
  - Chaum (1981): mix-nets.
  - Chaum (1991): group signatures.
  - Stajano and Anderson (1999): cocaine auction protocol.
- Applications: flexible, but more efficient in targeted communications. For example:
  - Voting systems.
  - Payments.

# *Vanilla mix-net*

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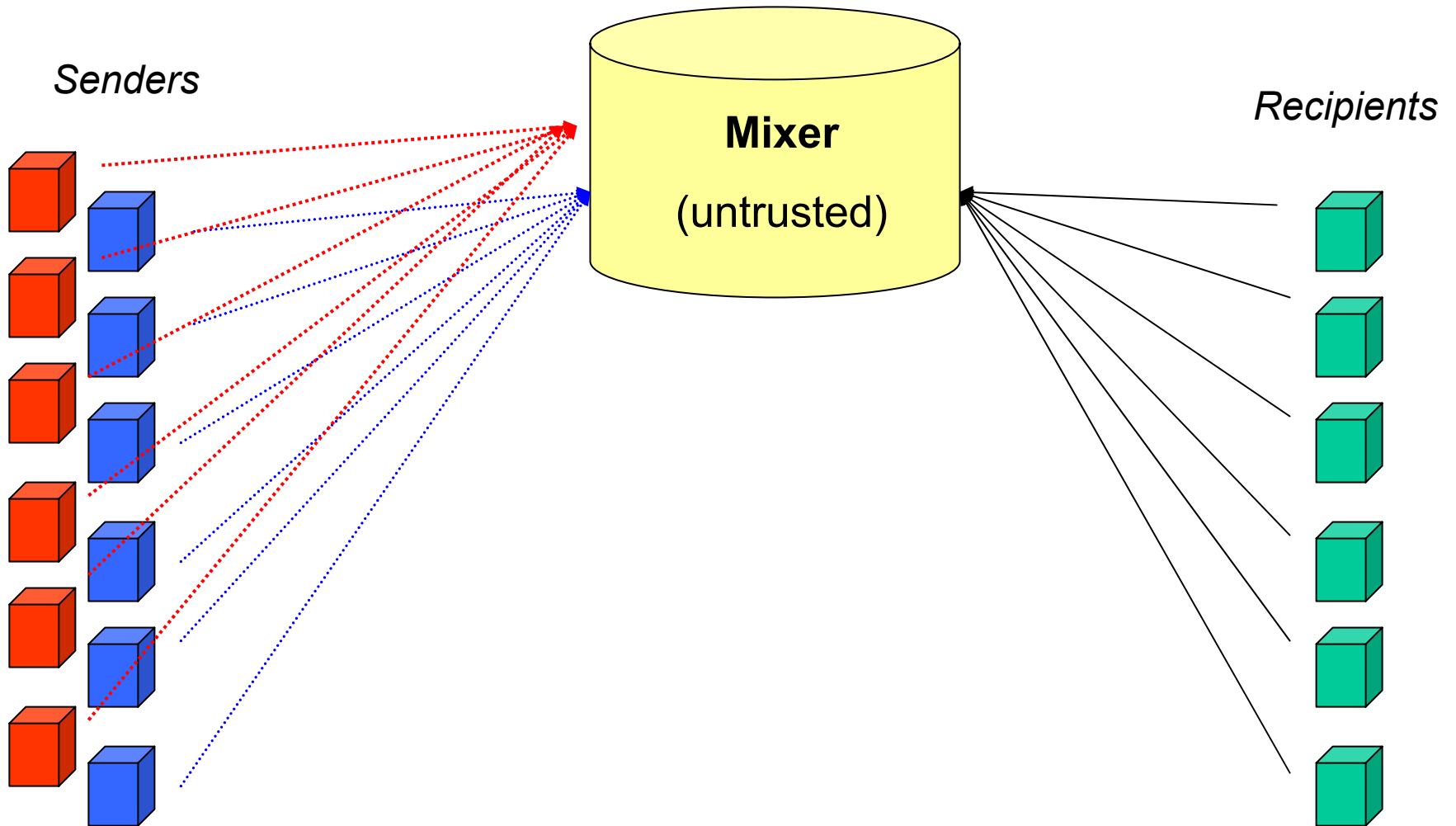
## *Issues discussed in the literature*

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- Trust.
- Reliability.
- Often, trade-offs between the two.

# *Vanilla acid mix*

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## *More precisely...*

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- Let users interact...
- ...through untrusted third party (mix)...
- ...splitting information...
- ...and broadcasting it.

# *Analysis*

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- Compare to Chaum (1981) voting mix-net protocol:
  - Candidate sends identification+key (pseudonym) through mix-net, then votes.
- Here:
  - Identification sent separately from key.
  - Mixed through other users.
- How?
  - Stajano and Anderson (1999). Message 3. can be broadcasted anonymously – does not contain identifying information (or, see Pfitzmann and Waidner [1986]).

# *Strengths, weaknesses, and attacks*

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- Strengths
  - Untrusted third party.
  - Untrusted senders.
  - Flexible.
- Weaknesses
  - Efficiency (depending on application).
- Attacks
  - Intersection attack.
  - Adversary observes in/out communication and owns *some* senders: OK.
  - Adversary sees in/out communication and owns *all* senders (“n-1 attack”): Not OK.

# *Applications*

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- (Messaging)
- Payments
  - Sender/buyer unlinkability.
- Voting
  - Receipt free.
  - Universally verifiable.
  - Open-ended ballot question.
  - (caveats.)

# *For the record*

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1.  $C_{C\pm} \rightarrow F : E_{C\pm PR} \{ E_{F\_PB} \{ C\_transaction\_id, C\_amount, C\_tPB, C\_tPBT, C\_t \} \}$
2.  $F \rightarrow C_{C\pm} : E_{C\pm PB} \{ C\_transaction\_id, T_{1,...,n}^{C\pm} \}$
3.  $C_{C\pm} \rightarrow * : E_{C\pm PR} \{ C\_tPB, n_{C\pm} \}$
4.  $[1, 2, ..., X]_{[1,2,...,X]_{-(t+1)}} \rightarrow C_{C\pm} :$   
 $E_{C\pm PB} \{ E_{[1,2,...,X]_{-(t+1)}PR} \{ [1, 2, ..., X]_{-(t+1)}PB, E_{C\pm PR} \{ C\_tPB, n_{C\pm}, \} \} \}$
5.  $C_{C\pm} \rightarrow F, [1, 2, ..., X]_{[1,2,...,X]_{-(t+1)}} :$   
 $E_{C\pm PR} \{ n_{C\pm}, [C, 1, 2, ..., X]_{-(t+1)}PB, [C, 1, 2, ..., X]_{-(t+1)}PBT \}$
6.  $[C, 1, 2, ..., X]_{[C,1,2,...,X]_{\pm}} \rightarrow F, [C, 1, 2, ..., X]_{[C,1,2,...,X]_{-(t+1)}} :$   
 $E_{[C,1,2,...,X]_{\pm}PR} \left\{ S, E_{F\_PB} \left\{ T_{1,...,n_{C\pm}}^{[C,1,2,...,X]_{\pm}}, [C, 1, 2, ..., X]_{\pm}PRT \right\} \right\}$
7.  $F \rightarrow * : E_{F\_PR} \{ E_{C_{-(t+1)}PB} \left\{ T_{1,...,n_{C\pm}}^{C_{-(t+1)}} \right\}, E_{1_{-(t+1)}PB} \left\{ T_{1,...,n_{C\pm}}^{1_{-(t+1)}} \right\}, ..., E_{X_{-(t+1)}PB} \left\{ T_{1,...,n_{C\pm}}^{X_{-(t+1)}} \right\} \}$

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