

Privacy-Preserving Interdomain Routing at Internet Scale

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Disclaimer

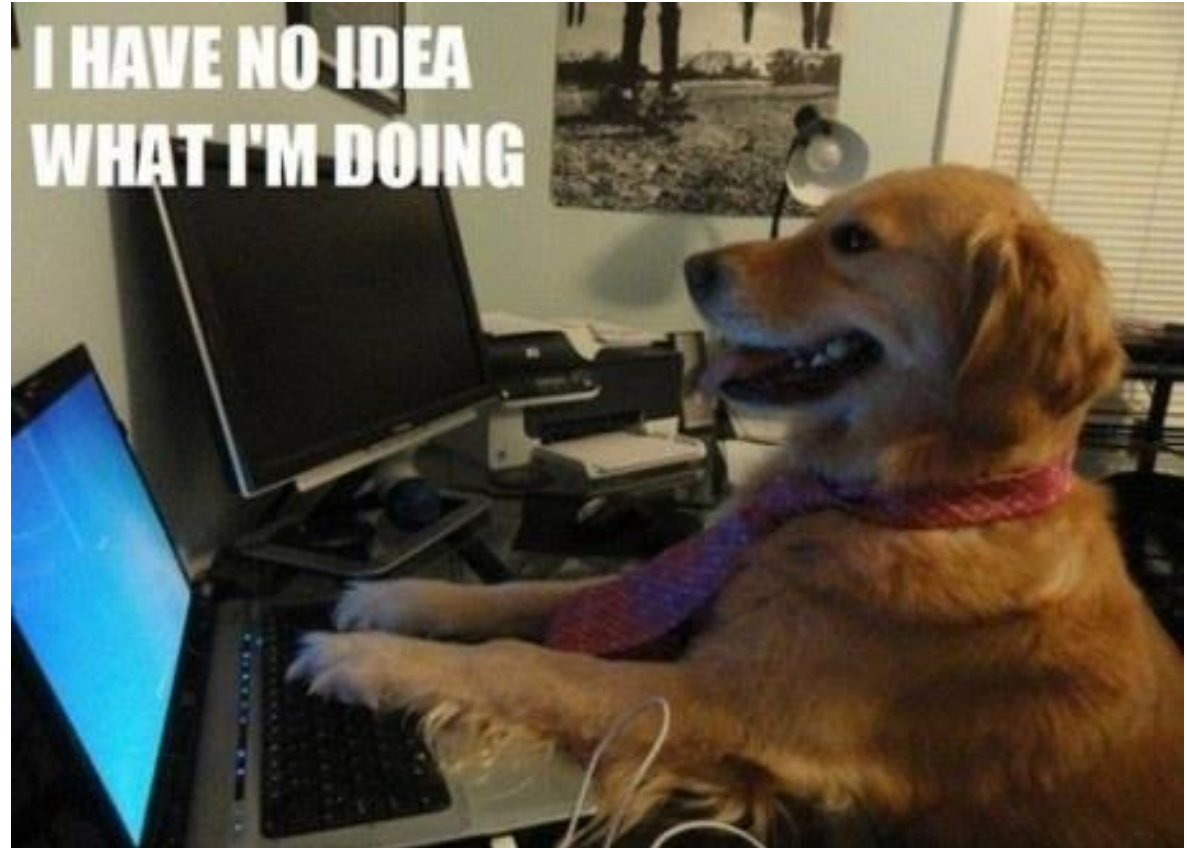


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Privacy-Preserving Interdomain Routing at Internet Scale

Disclaimer

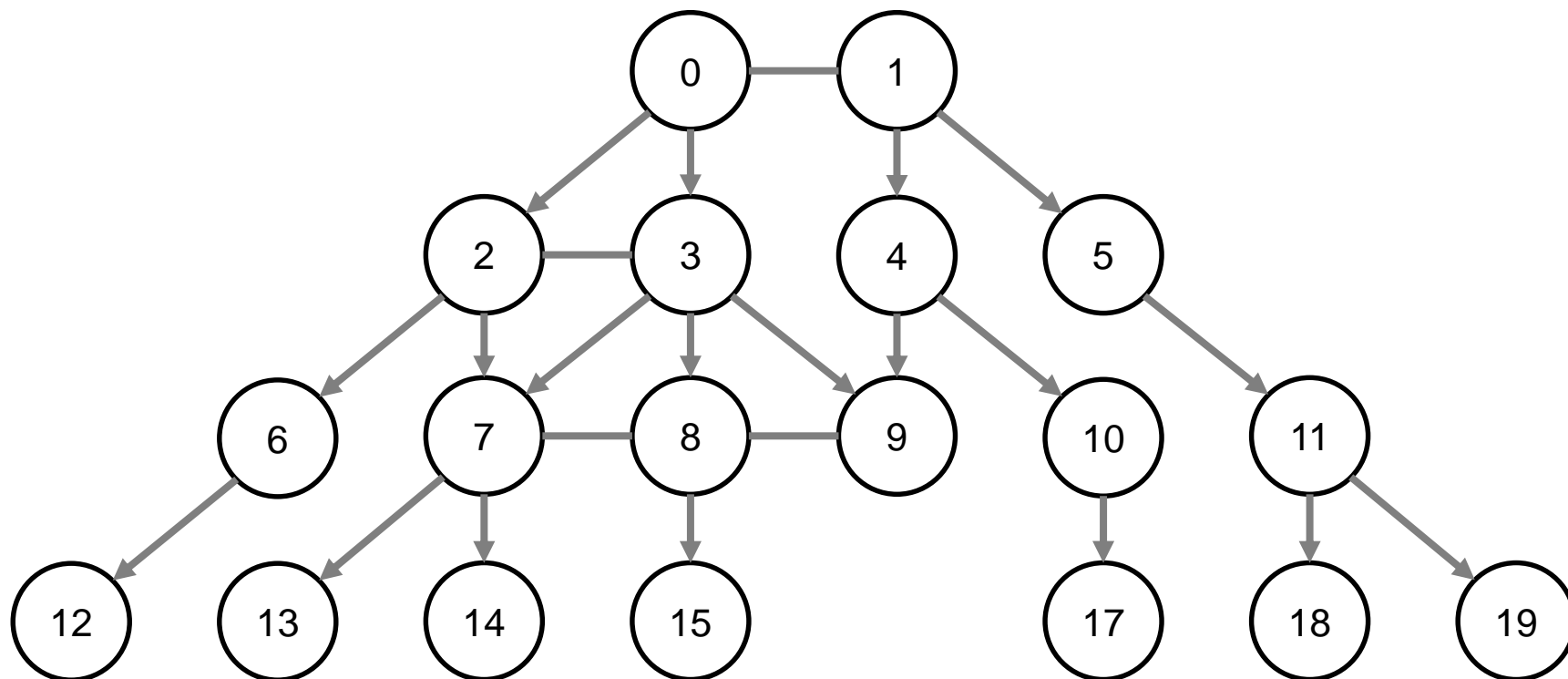
Privacy-Preserving Interdomain Routing at Internet Scale



What is BGP? (Approximation)



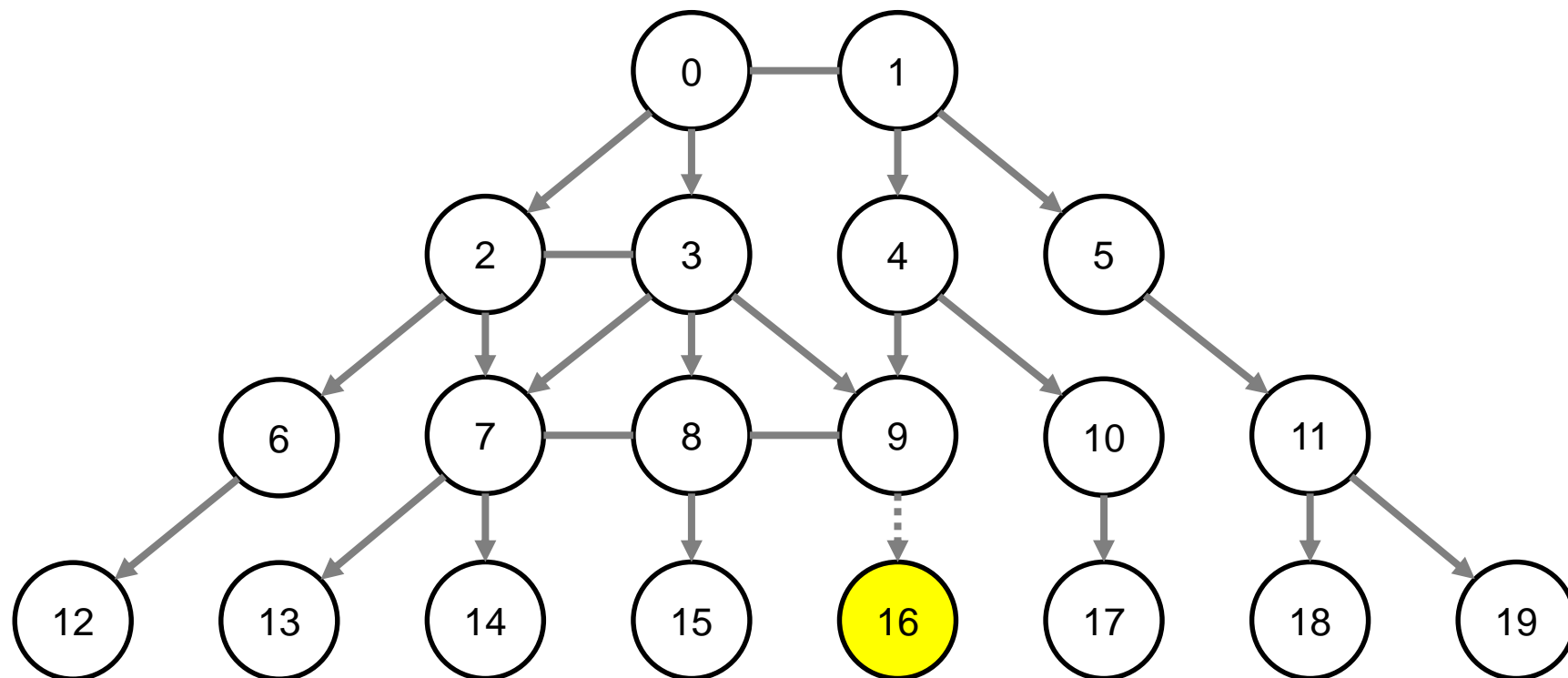
Compute all the routes to AS16.



What is BGP? (Approximation)



Compute all the routes to AS16.



What is BGP? (Approximation)



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Compute :

compute



Privacy-Preserving Inter-Domain Routing

Main problems of BGP: Convergence & Privacy

Original idea from [GSP⁺12] – Centralizing + SMPC!

Problem: only for toy example, impractical runtime.

This work:

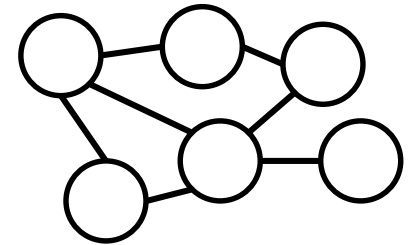
Real-world application of secure computation

56k autonomous systems with 239k connections!

We have two solutions that protect:

the **relations** between nodes: customer / provider or peering

the **export policy** and **preferences** of nodes



Generic Secure Two-Party Computation

First Ideas date back to 1980s

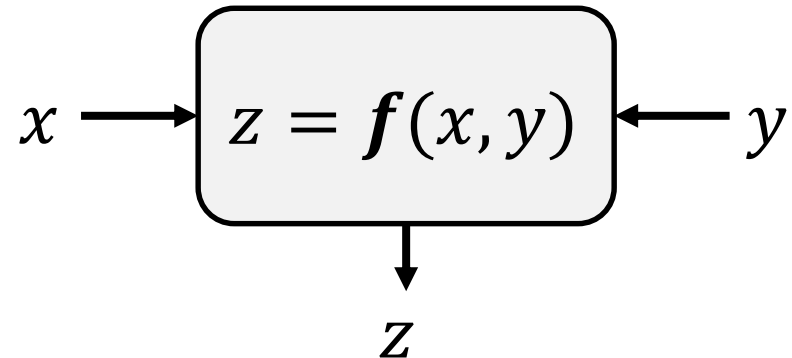
Generic applications



This work:

Two parties

Security against semi-honest
(passive) adversaries

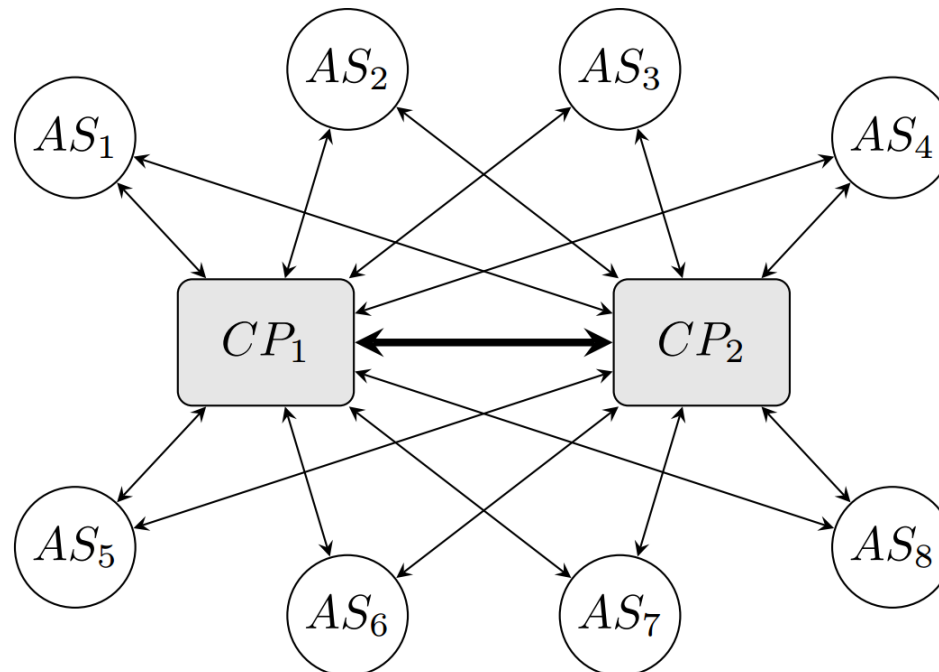


Privacy-Preserving Inter-Domain Routing

Centralized approach: Privacy-Issues solved by SMPC

2 computational parties (CPs), running our protocol

Each node (AS) secret-shares his private inputs with the CPs



Relation-Based Routing

Routing based on **relationship** between nodes:

Customers pay providers to route traffic

Peers route traffic for free

“*Economically driven*” routing instead of shortest paths

High-level Neighbor Relation Algorithm:

Plaintext input: **Topology, Target AS** – Private input: **EP - Relations**

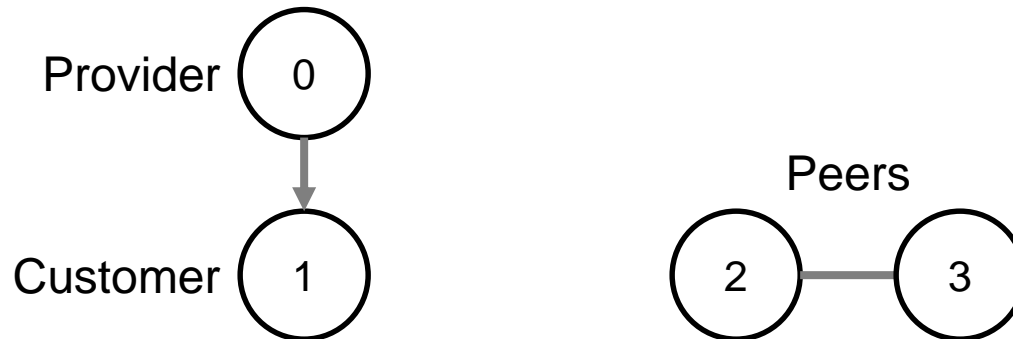
10 iterations for customer relation hops

1 iteration for peer hops

10 iterations for provider hops

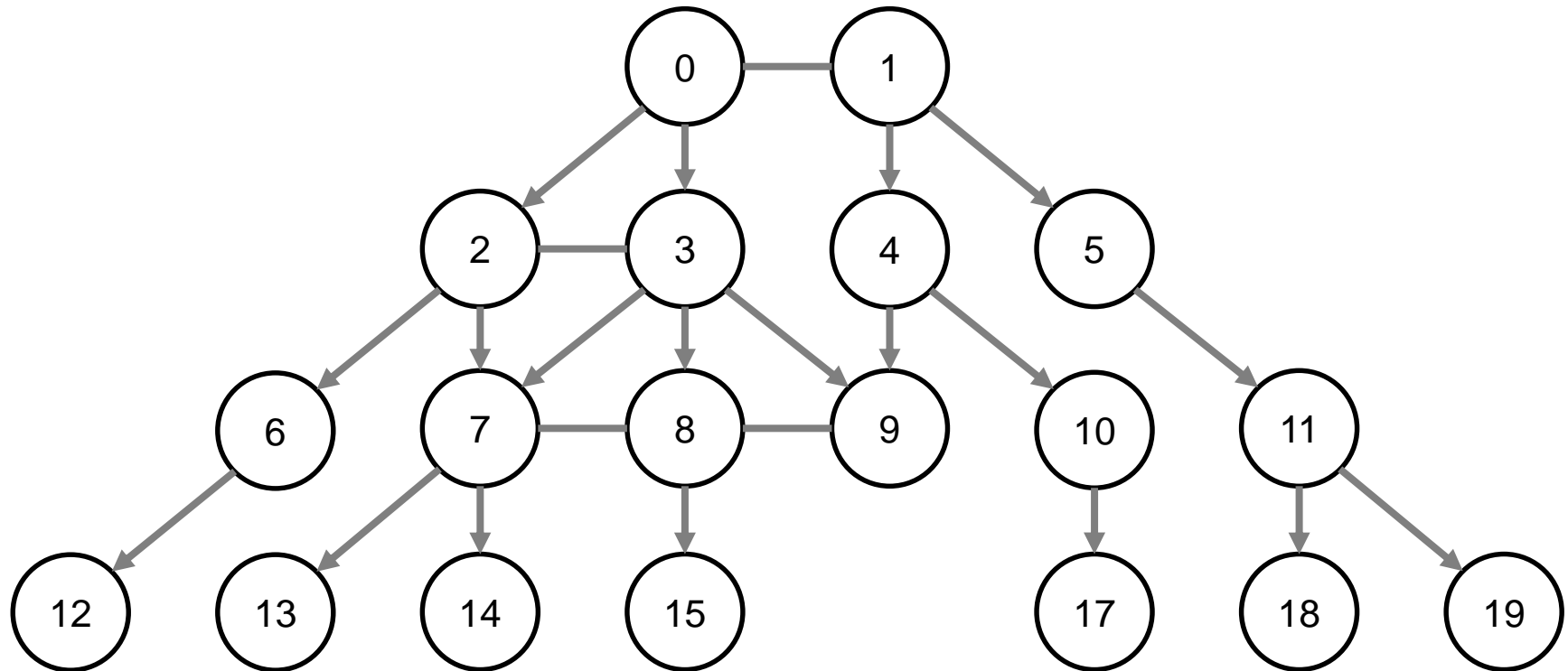
Private output: for every AS next hop to target AS

BGP Example – Notation



BGP Example

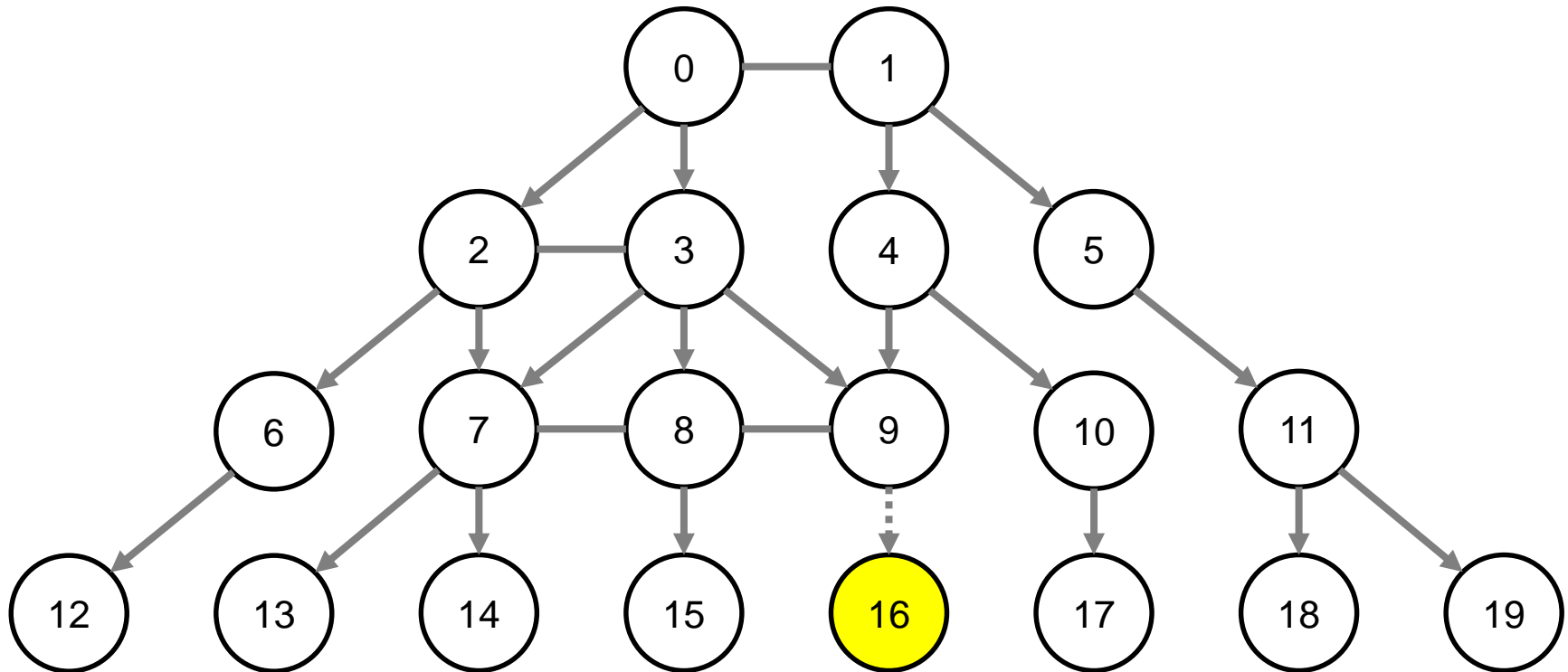
Public network topology



BGP Example

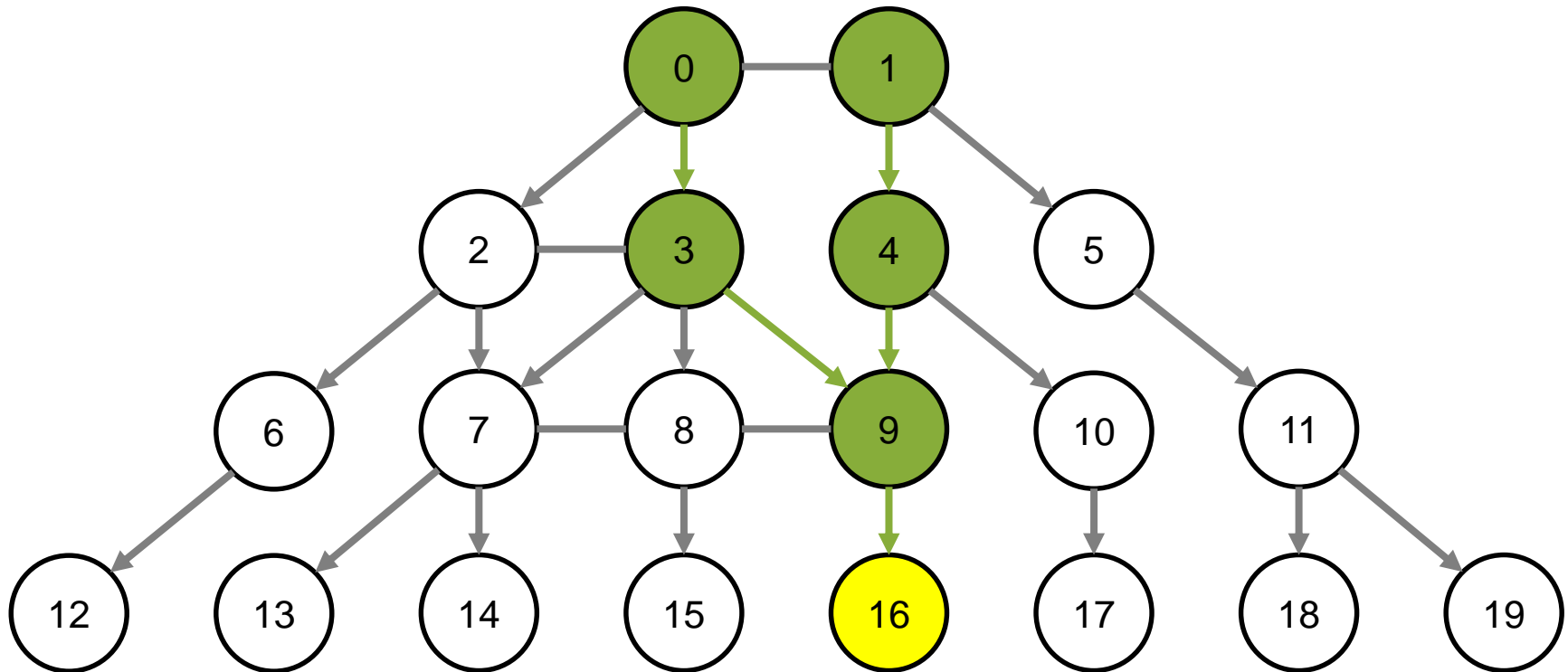
Public network topology

Node 16 is added



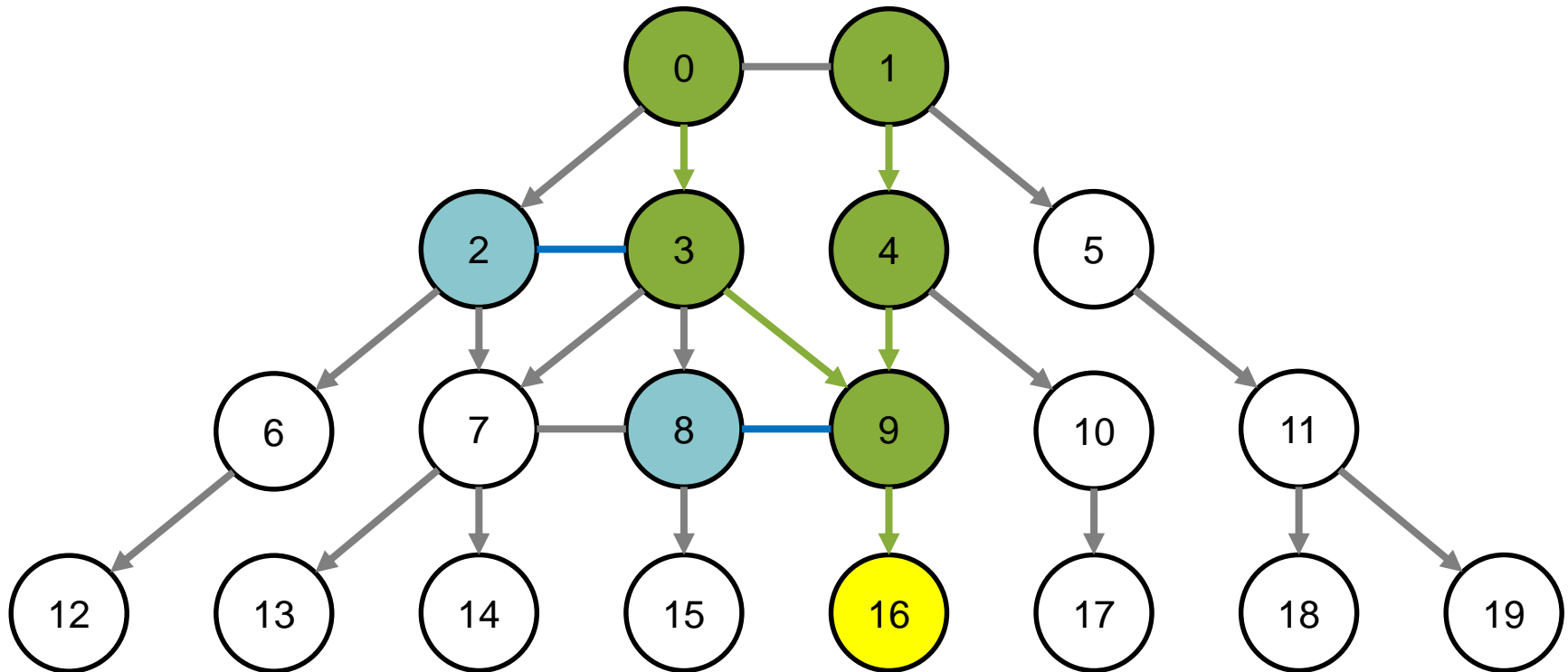
BGP Example

Routes through **customers** to 16



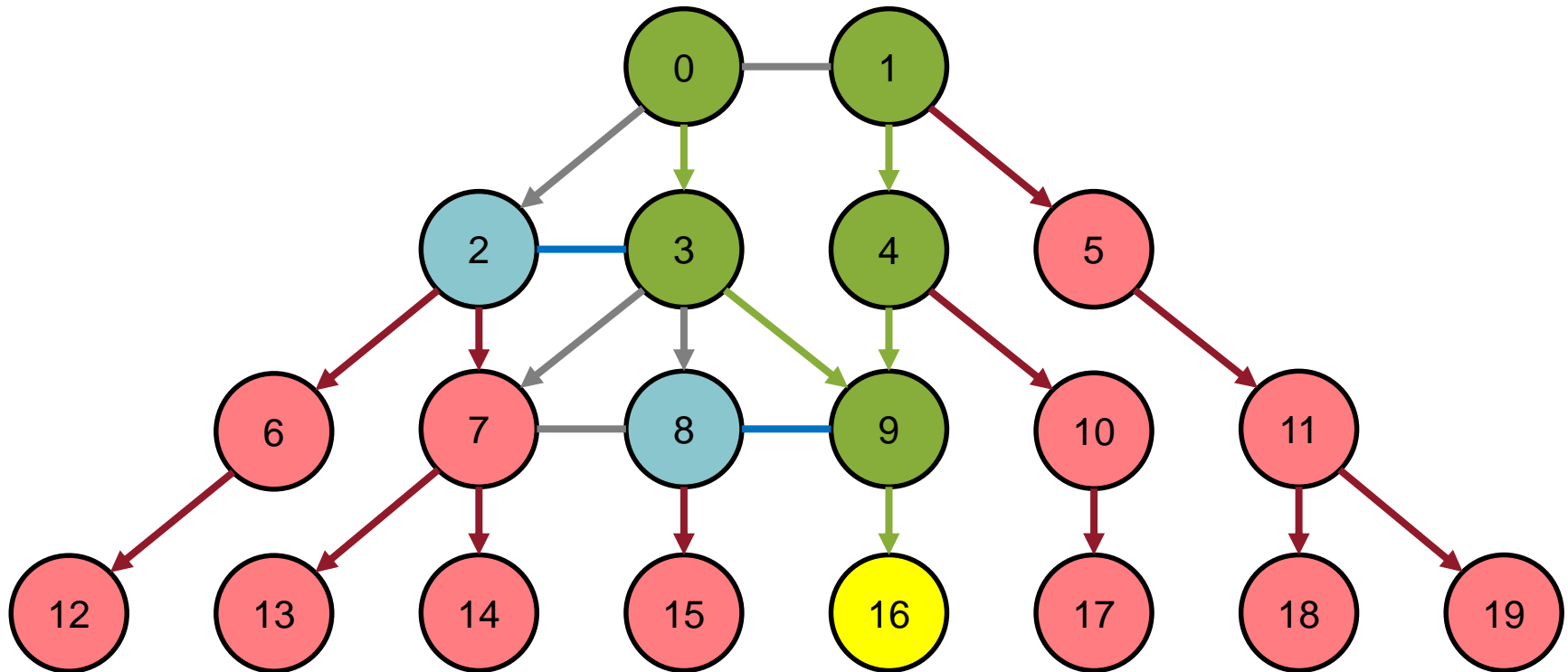
BGP Example

Routes through **peers** to 16



BGP Example

Routes through **providers** to 16



Preference-Based Routing



Routing based on **export policy** and **preference** between nodes:

ASes decide which routes are *published (exported)*

ASes have *preferences* for their neighbors

High-level Neighbor Preference Algorithm:

Plaintext input: **Topology, Target AS** – Private input: **EP - Preferences**

21 Iterations:

for all ASes:

for all of the ASes neighbors:

find highest **preference** neighbor with **published** route to **target**

Private output: for every AS next hop to target AS

Privacy-Preserving BGP – Circuit



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Algorithms built as **Boolean Circuit**:

SIMD operations

1 operation for multiple bits in parallel

Process all nodes in parallel on circuit level

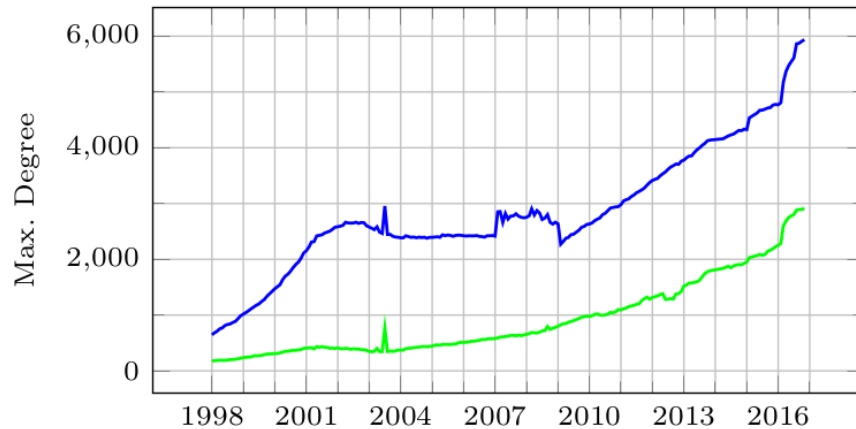
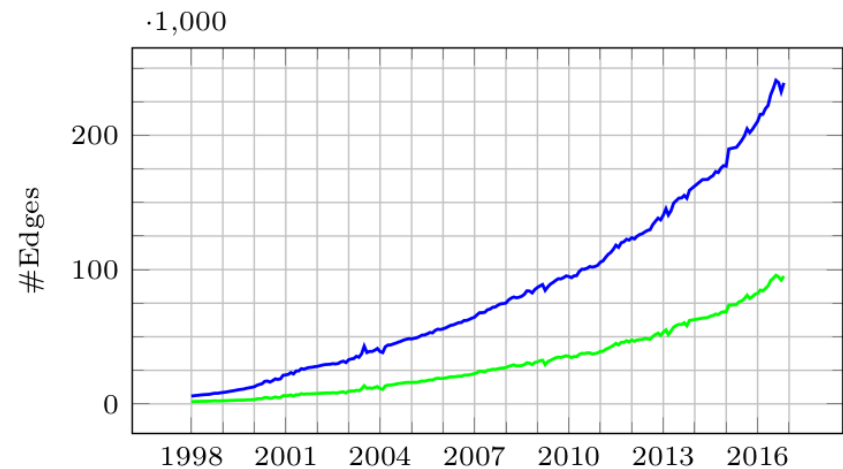
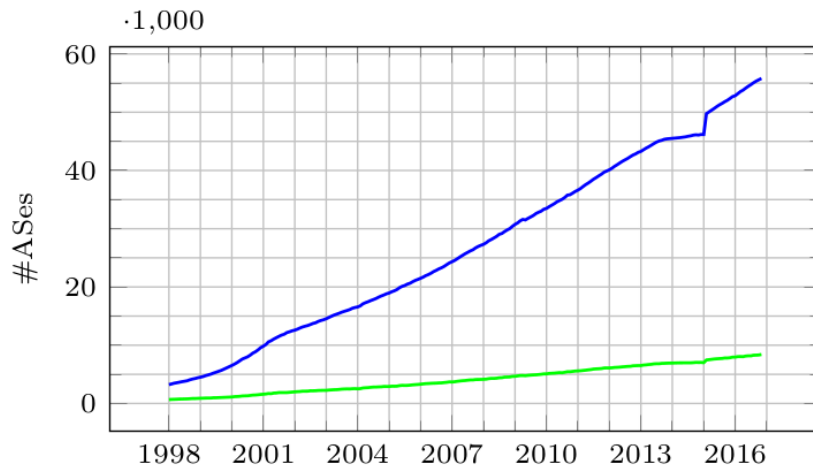
Efficient MUX with vector-ANDs in GMW

only 1 OT for n -bit values

Tree structure for depth-efficient parallel evaluation

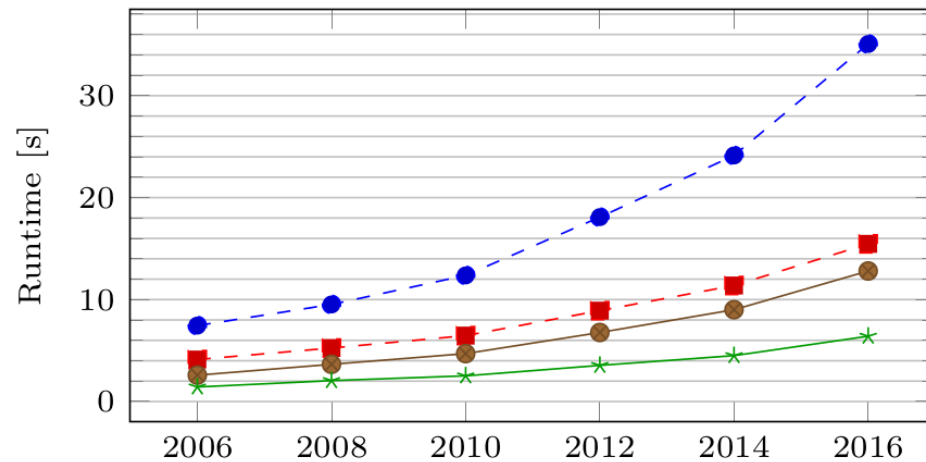
Algorithmic optimization: ignore stub nodes (85% of ASes)

CAIDA: BGP Statistics 1998 – 2016

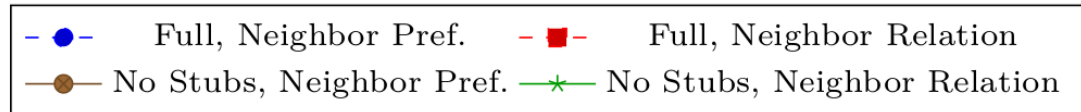
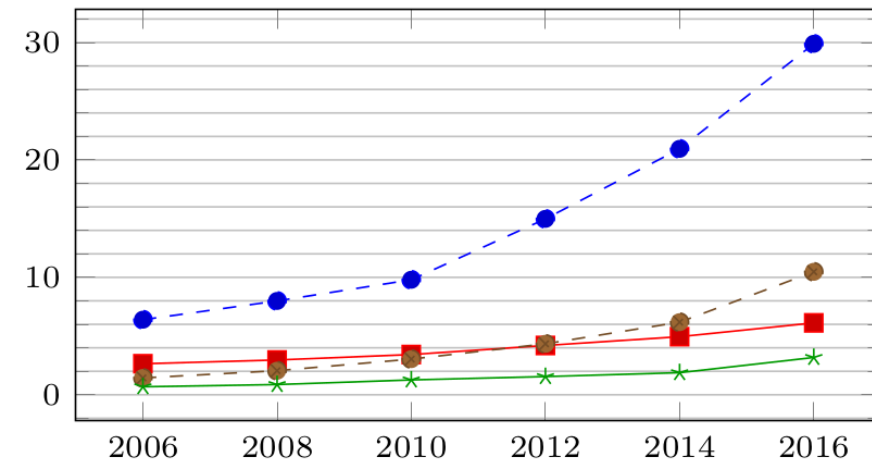


BGP Benchmarks: Full Topology

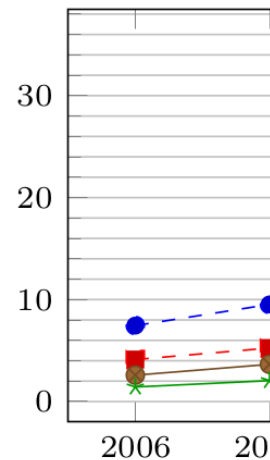
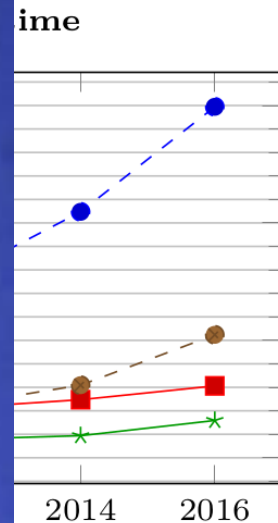
Setup Phase Runtime



Online Phase Runtime

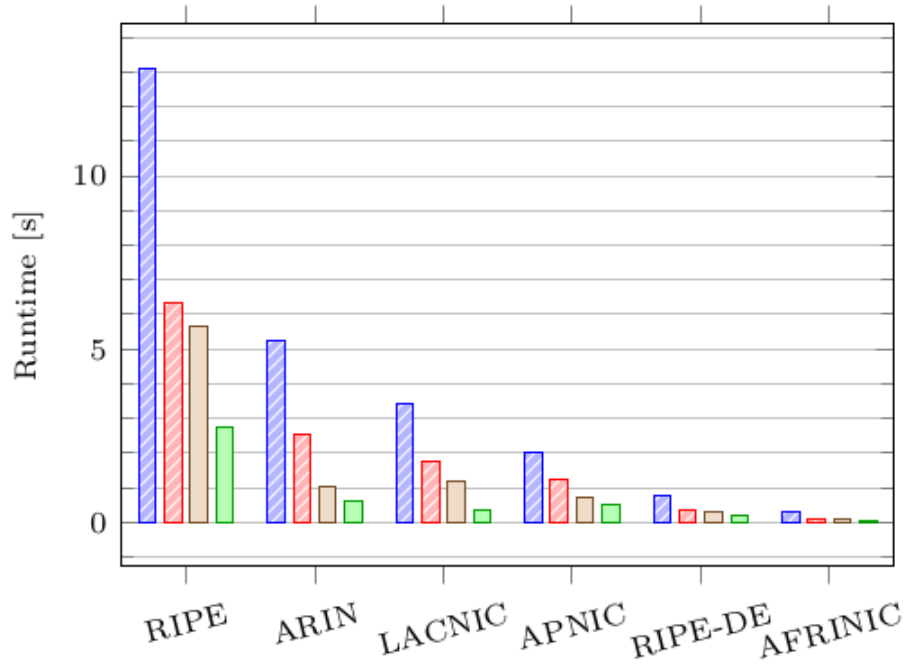


BGP Benchmarks: Full Topology

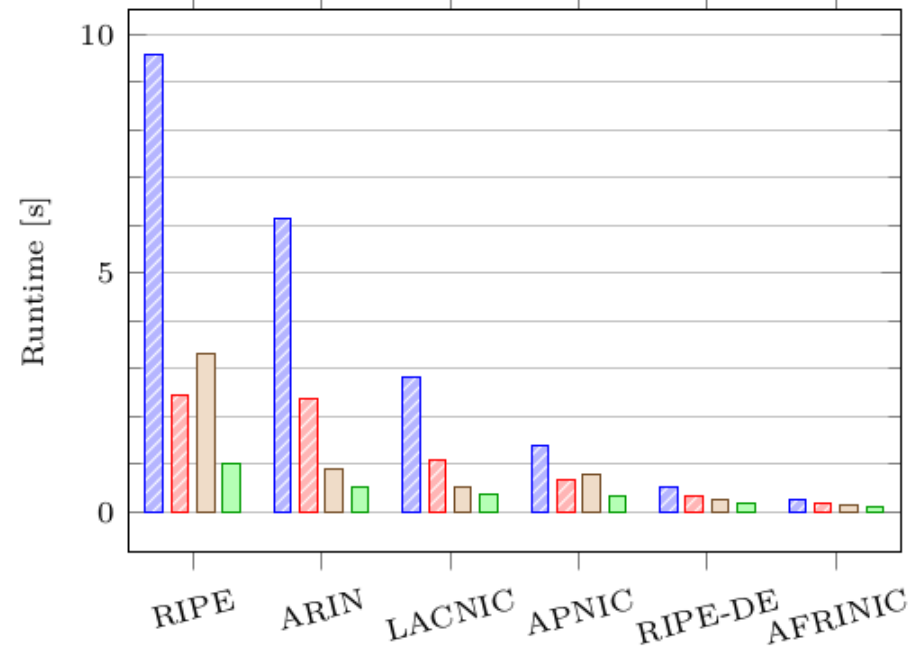


BGP Benchmarks: RIR Topology

Setup Phase Runtime



Online Phase Runtime



Legend: Full, Neighbor Pref. (blue hatched), Full, Neighbor Relation (red hatched), No Stubs, Neighbor Pref. (tan), No Stubs, Neighbor Relation (green)

Future Work and Conclusion

Hiding the topology?

Actual deployment?

Summary:

Real-World SMPC application

Made possible by algorithmic improvements and engineering



Thanks for your attention!

Questions?

References



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[GSP+12] D. Gupta, A. Segal, A. Panda, G. Segev, M. Schapira, J. Feigenbaum, J. Rexford, and S. Shenker. A new approach to interdomain routing based on secure multi-party computation. In *ACM Workshop on Hot Topics in Networks (HotNets'12)*, pages 37–42. ACM, 2012

Icons: <http://www.iconsmind.com>

ABY – A Framework for Efficient Mixed-Protocol Secure Two-Party Computation

C++ Framework for mixed-protocol secure two-party computation

Published at Network & Distributed Systems Symposium (NDSS'15)

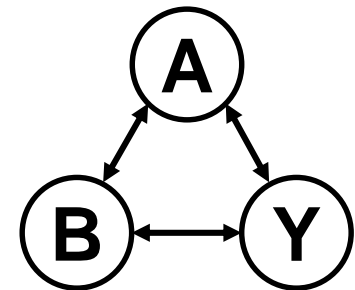
Multiple Protocols:

Arithmetic Sharing

Boolean Sharing (with the GMW protocol)

Yao's Garbled Circuits

Protocols split in **Setup** and **Online** phase



<http://www.encrypto.de/code/ABY>

