1. REVISITING PROTOCOL FLEXIBILITY

Internet technologies have been designed from guidelines like the robustness principle also known as Postel’s law \[1\]. Jon Postel’s law is described as: “Be conservative in what you do, be liberal in what you accept from others.” Fundamentally, it advises protocol designs to be tolerant with what they accept from the other peers. In practice, this law enables protocols to be forward compatible, meaning that they will be tolerant with a future version of the same protocol, making it possible to avoid unrecoverable errors when peers do not use the same protocol version. The Tor routing protocol naturally implements the robustness principle when processing data and control information, and it shows to be of crucial importance within a distributed and volunteer-based network that can be composed of nodes running many different versions of the protocol.

The robustness principle is elegant and straightforward to turn into a practical implementation. However, its security implications are often ignored. Rochet and Pereira \[2\] showed that the robustness principle could be exploited to convey information between malicious relays, find the guard relay used by a particular onion service or apply other various attacks. Furthermore, real attacks were observed in the history of the Tor network with techniques exploiting protocol robustness \[2\], raising awareness over the role this principle can have in efficiently implementing well-known theoretical attacks, like end-to-end traffic confirmation.

We propose to take a step back and wonder how the robustness principle could be revisited to support security requirements. Our goal would be to define a software architecture that offers the benefits of the robustness principle (i.e., efficient network services despite the presence of various software versions), while also guaranteeing that this robustness cannot be exploited by making sure that it is only used to support authentic evolutions of the protocol specification. We start to give an overview of the software architecture and then, we describe two families of usage scenarios: 1) lightweight updates for distributed systems: flexibility without Postel’s law; 2) Custom Internet Privacy.

2. FLEXIBLE ANONYMOUS NETWORK

We call FAN, for Flexible Anonymous Network, an anonymous network architecture able to transparently change its behavior for one or many users without having to restart relays or perturbing other user connections while proceeding to add, remove or modify protocol features. A FAN is achieved using what we call “Protocol Plugins”, a novel technology which will be further described in Section 3. We now describe two scenarios of our ongoing research that are made possible with a FAN architecture.

Lightweight updates. The first research question we consider is how to use Protocol Plugins to build a flexible Tor network which does not blindly comply with the robustness principle. Answering this research question would put Postel’s law behind us and may solve Tor’s security issues related to it, as well as offering new perspectives for feature deployment. A first perspective would be to offer Tor developers a method to propagate lightweight updates in a snap of fingers to all Tor users (relays, onion services, and Tor clients). Lightweight updates would offer flexibility without blindly following Postel’s law, since any new feature can be added to the protocol through Protocol Plugins and later be part of the codebase once the relay operator bumps Tor’s version through regular packaging management but, still, any deviation from the baseline protocol specification would need to be properly authenticated. The plugins are written in the same high-level language as Tor; they can easily be merged inside the main codebase for people running the up-to-date Tor version, leaving developers only to workout merge flows to manage plugins with the main codebase during development. To manage, download, delete and verify plugins authenticity, lightweight updates could follow TUF \[5\]’s proposal, (the update framework), an already widely used framework for secure updates aimed to survive keys compromise.

Custom Internet privacy. An exciting direction for a FAN architecture applied to Tor would be to let Tor users and onion services to plug code in their own Tor circuits. This could offer the opportunity for each user to benefit from the right anonymity/performance trade-off during their transient usage of the Tor network. For example, a Tor user could decide to plug a specific padding scheme over its middle relay while visiting an onion service. On the relay, the protocol extension would be ephemeral and only valid for this specific connection. This research direction envisions a new type of architecture for a distributed anonymous communication network, where setting the communication protocol to use at a given time would be itself private, meaning that the attacker would have the additional difficulty to apprehend which type of privacy-preserving technique is used to deanonimize the protected data flow.

However, such remote code injection capability is also potentially dangerous for the network performance, security, stability and for users’ privacy. We designed \[3\] a plugin management system that bears similarity to Certifi-
cate Transparency which allows a way to distribute adequate proofs that the plug-in enforces some claimed security properties. According to this design, relays, and end-users would only accept plugins with valid proofs that the code is consistent with some properties that they require to be guaranteed.

3. PROTOCOL PLUGINS

We suggest redesigning anonymous communication implementations such as Tor to make them flexible through Protocol Plugins, which are pieces of code that are executed inside a userland VM (yet with comparable performance to native execution) in response to a particular event.

We design Protocol Plugins from high-level languages such as C or Rust and we target a bytecode representation of a conceptual machine using the LLVM compiler Clang. In our experiments, we looked into eBPF’s virtual machine and WebAssembly’s virtual machine. eBPF’s virtual machine has been designed to run BPF bytecode inside the Linux kernel, like mini-programs which can be started in response to a particular event inside the kernel. This BPF bytecode is JITted to machine code and offers comparable to native performance while executing dynamically plugged programs. On the other hand, WebAssembly has been designed to offer an alternative to JavaScript regarding performance on critical tasks and to offer a higher diversity of languages for the web Platform since .wasm bytecode can be obtained from the compilation of diverse high-level languages. Figure 1 offers a representative view of the process to build Protocol Plugins. Both BPF and wasm’s execution can be emulated during development, which makes debugging a lot more comfortable but still translates to machine code in production. WebAssembly in its bytecode format has nothing dependent to the Web, nor eBPF has dependencies to the Linux kernel. Both are only instructions in a low-level format that follows a given conceptual machine architecture.

![Diagram of Protocol plugin lifecycle](image)

**Figure 1: Protocol plugin lifecycle**

Protocol Plugins are, in our design, embedded in Tor’s critical path when handling data cells and would be called for any protocol operation that would not follow the native implementation. We distinguish between the operations handled by the native code which each relay runs, depending on their Tor version, and the operations handled by protocol plug-ins as default operations and non-default operations (what is default or non-default then depends on the specific Tor version installed on a node).

Non-default operations call Protocol Plugins deployed on relays in order to follow the up-to-date version of the Tor Routing Protocol. If no Protocol Plugin can handle the data, nor the default operation, it means that the information received does not match the protocol’s specifications. In consequences, the relays could apply the most efficient error-handling mechanism to protect users’ privacy, such as killing the Tor circuit and reporting the error if needed. Protocol Plugins are a few KBs files containing the bytecode produced by the LLVM compiler. Portable by design, Protocol Plugins can be distributed to relays independently from their system architecture and run inside optimized sandboxes with comparable to native performance once JITed to machine code.

We made a prototype implementation that is already working on the receiving side of plugins and allows plugging and executing JITed BPF bytecode in response to Tor’s protocol events inside a sandboxed userland VM. New protocol operations are plugged to Tor’s main code in less than 1ms on a regular laptop (8th Gen Intel, 2666Mhz DDR4 and NVMe SSD), which accounts for the load of the plugin from the disk and the bootstrap of the user space virtual machine used to sandbox the plugin. The current implementation of the virtual machine allows exporting functions from the host application (Tor in this case) at compile time (e.g., helper functions). For more dynamical needs, like accessing some internal Tor data from a plugin, we defined an interface between the host application and the plugins in order to allow the plugin to get/set Tor’s internal state since, by design, we cannot dereference pointers outside of the space allocated for the virtual machine. Data access capabilities can be given to the virtual machine running the plugin, including Tor’s internal data or even external access, like authorizing a system call to open a particular file (e.g., writing logs) but nothing else. Tor’s code calls the plugin in response to some particular event, like the reception of a cell that includes an unhandled protocol feature within its header. In that case, the cell would be transmitted to the plugin handling that particular protocol feature. Any other application-defined event can also call plugins.

4. A NEW PARADIGM?

Anonymous communications is one of many fields that could benefit from protocol plugins. Protocol plugins may serve other goals, such as advancing the arm race in censorship resistance (e.g., exploring how an authorized protocol or application could hide another protocol from the censor through plugins), advancing deployment’s speed of new transport protocols or extensions [3], advancing the Internet’s control plane interoperability, etc.

5. REFERENCES


